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<u>L18</u>	l17 and (flush\$4 or cache or modif\$4 or updat\$4 or tag or buffer or queue)	22	<u>L18</u>
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<u>L16</u>	L14 and l13	54	<u>L16</u>
<u>L15</u>	L14 and l13 and l12	5	<u>L15</u>
<u>L14</u>	(("L1" or upper or higher or first) adj2 cache) near8 (flush\$4 or swap\$4 or push\$4) near8 (("L2" or lower or second) adj2 cache)	63	<u>L14</u>
<u>L13</u>	cache near6 (modif\$8 or updat\$4 or dirty or (store adj in))	5284	<u>L13</u>
<u>L12</u>	(flush or swap) adj2 buffer	598	<u>L12</u>
<u>L11</u>	L10 and ((711/122 711/133 711/135)!.CCLS.)	35	<u>L11</u>
<u>L10</u>	L9 and l3	213	<u>L10</u>
<u>L9</u>	buffer near4 (flush\$3 or clear\$3 or clean\$3 or empt\$4 or remov\$4 or eliminat\$4) near8 cache	330	<u>L9</u>
<u>L8</u>	l5 and (buffer near4 (flush\$3 or clear\$3 or clean\$3 or empt\$4 or remov\$4 or eliminat\$4))	600	<u>L8</u>
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<u>L4</u>	cache near8 (flush\$3 or clear\$3 or clean\$3 or empt\$4 or remov\$4 or eliminat\$4)	4541	<u>L4</u>
<u>L3</u>	cache near8 (modif\$8 or updat\$4 or (store adj in))	5299	<u>L3</u>
<u>L2</u>	L1 and (cache near8 (modif\$8 or updat\$4 or (store adj in)))	8	<u>L2</u>
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Search Results

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- 1** **Classifying load and store instructions for memory renaming** 96%
Glenn Reinman , Brad Calder , Dean Tullsen , Gary Tyson , Todd Austin
Proceedings of the 13th international conference on Supercomputing
May 1999
- 2** **Memory dependence prediction using store sets** 84%
George Z. Chrysos , Joel S. Emer
ACM SIGARCH Computer Architecture News , Proceedings of the 25th annual international symposium on Computer architecture April 1998
Volume 26 Issue 3
For maximum performance, an out-of-order processor must issue load instructions as early as possible, while avoiding memory-order violations with prior store instructions that write to the same memory location. One approach is to use memory dependence prediction to identify the stores upon which a load depends, and communicate that information to the instruction scheduler. We designate the set of stores upon which each load has depended as the load's "store set". The processor can discover and u ...
- 3** **Memory forwarding: enabling aggressive layout optimizations by guaranteeing the safety of data relocation** 80%
Chi-Keung Luk , Todd C. Mowry
ACM SIGARCH Computer Architecture News , Proceedings of the 26th annual international symposium on Computer architecture May 1999

Volume 27 Issue 2

By optimizing data layout at run-time, we can potentially enhance the performance of caches by actively creating spatial locality, facilitating prefetching, and avoiding cache conflicts and false sharing. Unfortunately, it is extremely difficult to guarantee that such optimizations are *safe* in practice on today's machines, since accurately updating *all* pointers to an object requires perfect alias information, which is well beyond the scope of the compiler for languages such as C. T ...

4 Implementation of precise interrupts in pipelined processors 80%

James E. Smith , Andrew R. Pleszkun

25 years of the international symposia on Computer architecture
(selected papers) August 1998

5 Cache replacement with dynamic exclusion 80%

Scott McFarling

ACM SIGARCH Computer Architecture News , Proceedings of the 19th
annual international symposium on Computer architecture April 1992
Volume 20 Issue 2

Most recent cache designs use direct-mapped caches to provide the
fast access time required by modern high speed CPU's.

Unfortunately, direct-mapped caches have higher miss rates than
set-associative caches, largely because direct-mapped caches are
more sensitive to conflicts between items needed frequently in the
same phase of program execution. This paper presents a new
technique for reducing direct-mapped cache misses caused by
conflicts for a particular cache line. A small fi ...

6 Instruction fetch energy reduction using loop caches for
embedded applications with small tight loops 77%

Lea Hwang Lee , Bill Moyer , John Arends

Proceedings 1999 international symposium on Low power electronics
and design August 1999

7 Tolerating late memory traps in ILP processors 77%

Xiaogang Qiu , Michel Dubois

ACM SIGARCH Computer Architecture News , Proceedings of the 26th
annual international symposium on Computer architecture May 1999
Volume 27 Issue 2

ILP processors can execute a large number of instructions at the
same time. Thus it becomes more and more difficult to support
traps efficiently. On the other hand a current trend in architecture
is to support various memory functions in software rather than
hardware, usually by trapping the execution processor on a cache

miss, TLB miss or a failed access to a local or remote memory. These late memory traps block the faulting instruction at the top of the active list, backing up the pipeline. Mo ...

- 8** A novel renaming scheme to exploit value temporal locality 77%
through physical register reuse and unification
Stephen Jourdan , Ronny Ronen , Michael Bekerman , Bishara Shomar , Adi Yoaz
Proceedings of the 31st annual ACM/IEEE international symposium on Microarchitecture November 1998
- 9** Active memory: a new abstraction for memory system 77%
simulation
Alvin R. Lebeck , David A. Wood
ACM Transactions on Modeling and Computer Simulation (TOMACS) January 1997
Volume 7 Issue 1
- 10** Don't use the page number, but a pointer to it 77%
André Seznec
ACM SIGARCH Computer Architecture News , Proceedings of the 23rd annual international symposium on Computer architecture May 1996
Volume 24 Issue 2
Most newly announced high performance microprocessors support 64-bit virtual addresses and the width of physical addresses is also growing. As a result, the size of the address tags in the L1 cache is increasing. The impact of on chip area is particularly dramatic when small block sizes are used. At the same time, the performance of high performance microprocessors depends more and more on the accuracy of branch prediction and for reasons similar to those in the case of caches the size of the Br ...
- 11** Recovery protocols for shared memory database systems 77%
Lory D. Molesky , Krithi Ramamritham
ACM SIGMOD Record , Proceedings of the 1995 ACM SIGMOD international conference on Management of data May 1995
Volume 36 Issue 7
- 12** On reconfigurable on-chip data caches 77%
Fredrik Dahlgren , Per Stenström
Proceedings of the 24th annual international symposium on Microarchitecture September 1991
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